

Why Distributed Optimization?

Optimization in Networkland: Challenges and Opportunities of Distributed Methods

Giuseppe Notarstefano

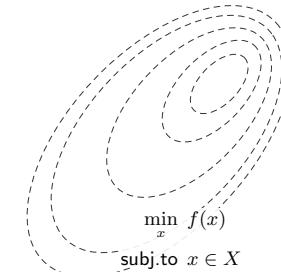
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June 28, 2019, Naples (Italy)



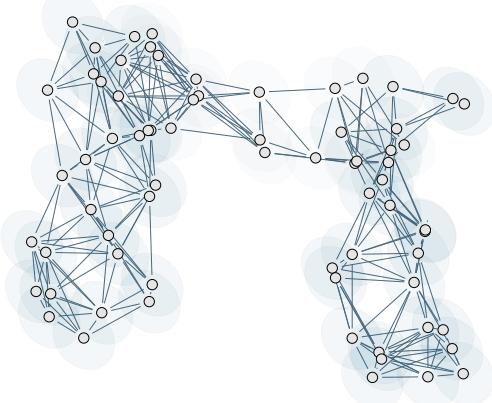
Mathematical problems appearing in several fields
 (Engineering, Economics, Biology, Social Sciences, ...)

Well-established numerical schemes



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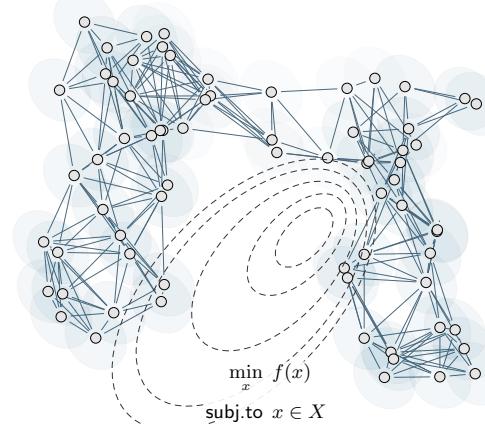
Why Distributed Optimization?



- Massive computation and communication
 - Local private data
 - Distributed algorithms (e.g., average consensus)
- ↓
- Share computation instead of data

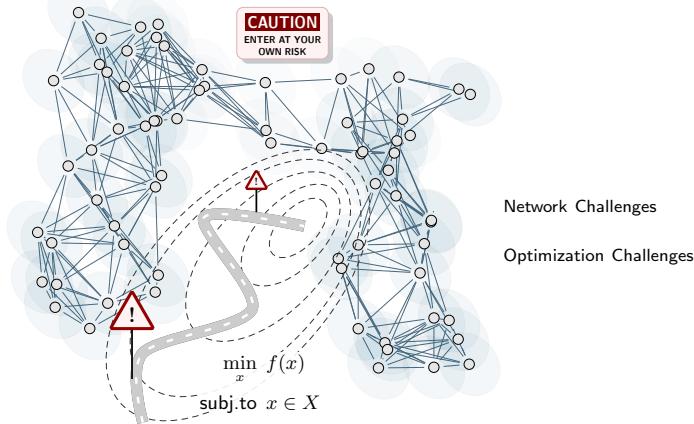
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Why Distributed Optimization?



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Network Challenges

Distributed algorithm: compute locally & communicate with neighbors

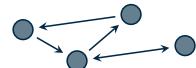


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Network Challenges

Distributed algorithm: compute locally & communicate with neighbors

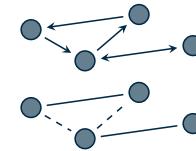
- Undirected/Directed



Network Challenges

Distributed algorithm: compute locally & communicate with neighbors

- Undirected/Directed
- Fixed/Time-Varying



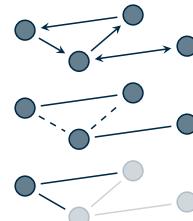
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Network Challenges

Distributed algorithm: compute locally & communicate with neighbors

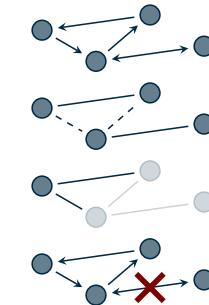
- Undirected/Directed
- Fixed/Time-Varying
- Asynchronous



Network Challenges

Distributed algorithm: compute locally & communicate with neighbors

- Undirected/Directed
- Fixed/Time-Varying
- Asynchronous
- Unreliable



Topology and communication NOT a design parameter

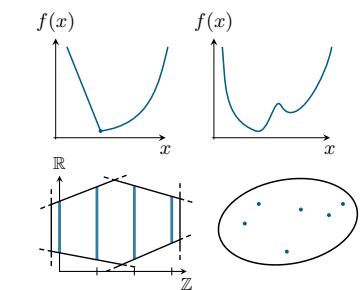
Optimization Challenges

$$\begin{aligned} \min_x f(x) \\ \text{subj.to } x \in X \end{aligned}$$

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Nonsmooth (convex), Nonconvex,
Mixed-Integer, Combinatorial



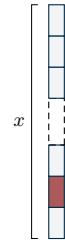
Optimization Challenges



$$\min_x f(x)$$

Nonsmooth (convex), Nonconvex,
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Big-Data (high dim. dec. var.)



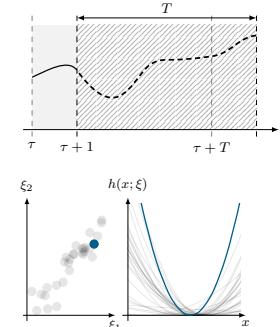
Optimization Challenges

$$\begin{array}{ll} \min\limits_x & f(x) \\ \text{subj.to} & x \in X \end{array}$$

Nonsmooth (convex), Nonconvex Mixed-Integer, Combinatorial

Big-Data (high dim. dec. var.)

Dynamic, Online,
Stochastic, Uncertain



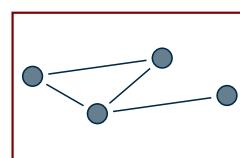
Distributed Optimization Paradigm



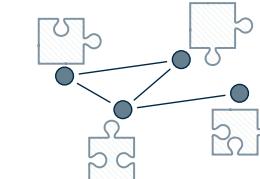
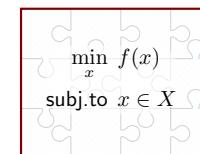
Optimization

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Network



Distributed Optimization Paradigm



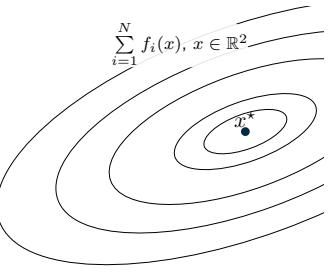
Network

- know only part of optimization problem
 - cooperate to compute a solution

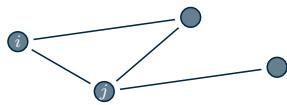
Cost-Coupled Set-up

$$\min_{x \in X} \sum_{i=1}^N f_i(x)$$

OPT4SMART



- N agents communicate over graph \mathcal{G}

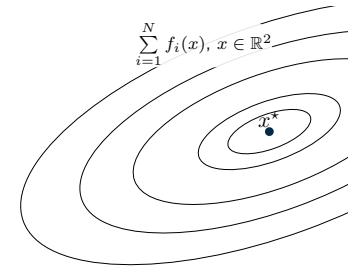


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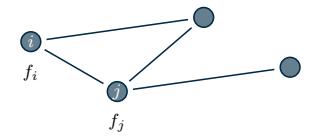
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- agent i knows f_i only

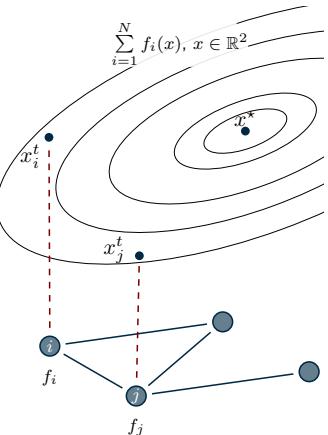


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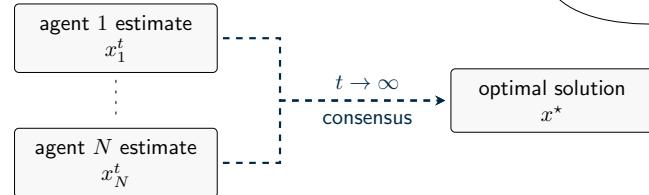
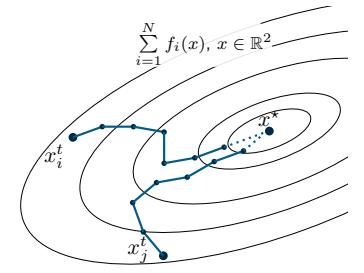
- N agents communicate over graph \mathcal{G}
- agent i knows f_i only
- x_i^t solution estimate of i

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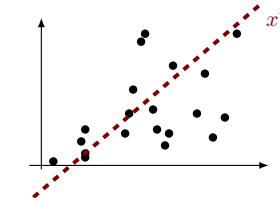
Data Analytics

$$\min_x \sum_{i=1}^N \underbrace{\|b_i - D_i x\|^2}_{f_i(x)} + r(x)$$



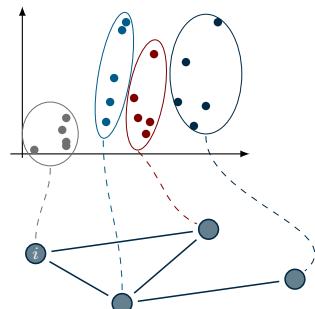
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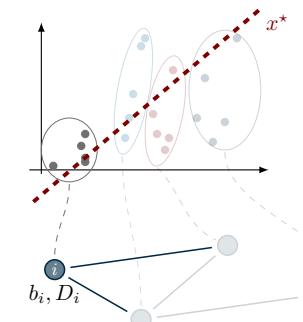


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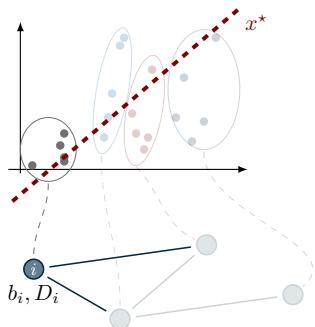
Paradigm

- local private data
- cooperate to learn from all data



Data Analytics

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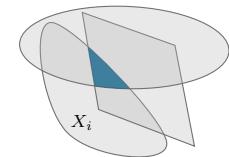


Paradigm

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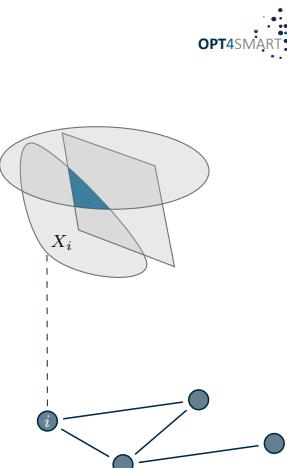
Share computation instead of data

$$\begin{aligned} & \min_x f(x) \\ \text{subj.to } & x \in \bigcap_{i=1}^N X_i \end{aligned}$$



Common Cost Set-up

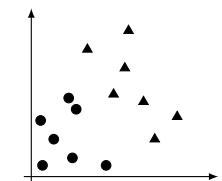
$$\begin{aligned} & \min_x f(x) \\ \text{subj.to } & x \in \bigcap_{i=1}^N X_i \end{aligned}$$



- agent i knows common cost f and local constraint X_i
- algorithms suited for this set-up

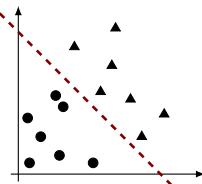
Support Vector Machine

$$\begin{aligned} & \min_{w,b,\xi} \frac{1}{2} \|w\|^2 + C\xi \\ \text{subj.to } & \ell_i(w^\top p_i + b) \geq 1 - \xi, \forall i \\ & \xi \geq 0 \end{aligned}$$



Support Vector Machine

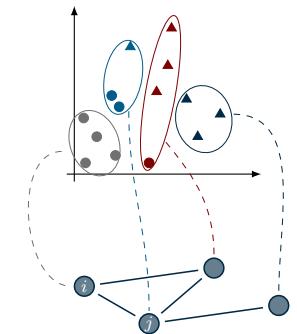
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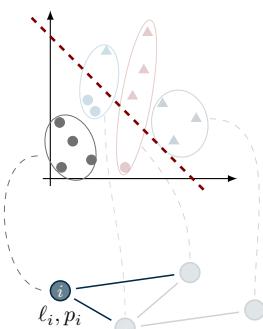
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Local constraints and common cost



Support Vector Machine

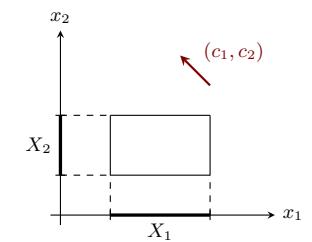
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Local constraints and common cost

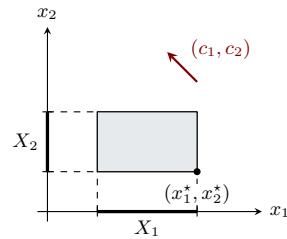
Constraint-Coupled Set-up

$$\begin{aligned} & \min_{x_1,x_2} c_1 x_1 + c_2 x_2 \\ \text{subj.to } & x_1 \in X_1, x_2 \in X_2 \end{aligned}$$



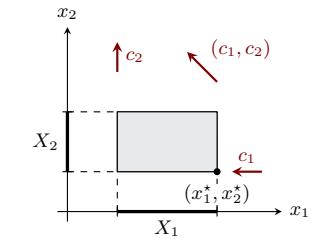
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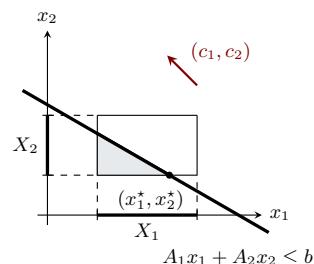
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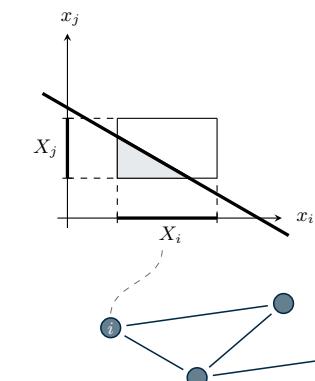
Constraint-Coupled Set-up

$$\begin{aligned} \min_{x_1, x_2} \quad & c_1 x_1 + c_2 x_2 \\ \text{subj.to} \quad & x_1 \in X_1, x_2 \in X_2 \\ & A_1 x_1 + A_2 x_2 \leq b \end{aligned}$$



Constraint-Coupled Set-up

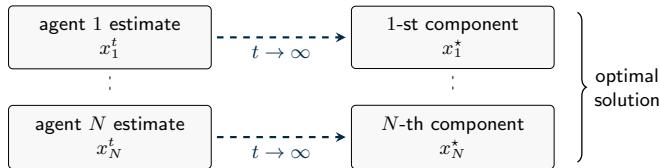
$$\begin{aligned} \min_{x_1, \dots, x_N} \quad & \sum_{i=1}^N f_i(x_i) \\ \text{subj.to} \quad & x_i \in X_i, \forall i \\ & \sum_{i=1}^N g_i(x_i) \leq 0 \end{aligned}$$



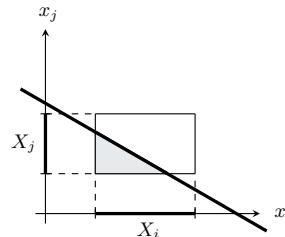
- \$N\$ agents communicate over graph \$\mathcal{G}\$
- \$(x_1, \dots, x_N)\$ dec. var. – size grows with \$N\$
- agent \$i\$ knows \$f_i, g_i\$ and \$X_i\$

Constraint-Coupled Set-up

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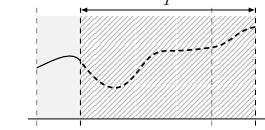
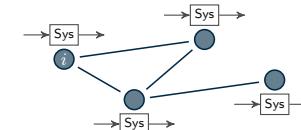


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Distributed Optimal Control

$$\min_{z_1, \dots, z_N} \quad \sum_{i=1}^N \left(\sum_{\tau=0}^T \ell_i(z_i(\tau), u_i(\tau)) + m_i(z_i(T)) \right)$$



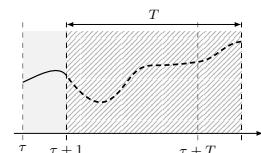
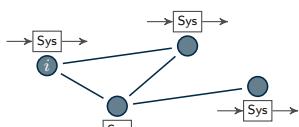
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$$\text{subj.to } z_i(\tau+1) = A_i z_i(\tau) + B_i u_i(\tau), \forall i, \tau \in [0, T]$$

$$z_i(\tau) \in Z_i, \quad u_i(\tau) \in U_i, \quad \forall i, \tau \in [0, T]$$



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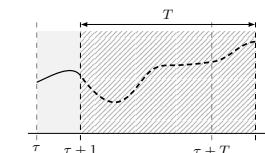
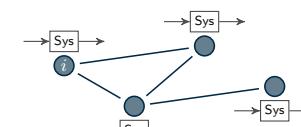
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$$\sum_{i=1}^N H_i z_i(\tau) \leq h, \quad \tau \in [0, T]$$



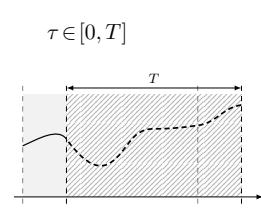
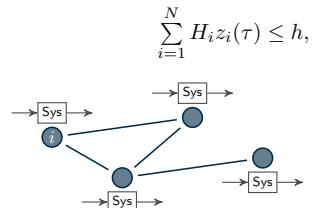
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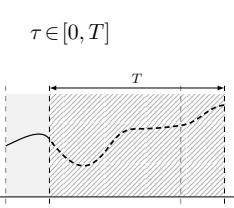
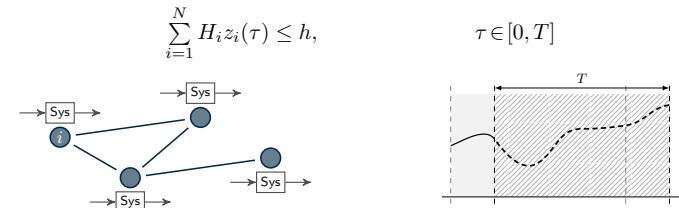
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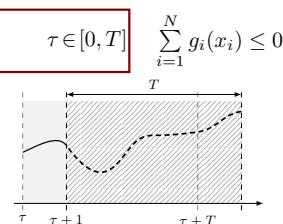
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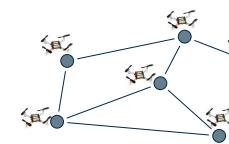
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$$\text{subj.to } z_i(\tau+1) = A_i z_i(\tau) + B_i u_i(\tau), \forall i, \tau \in [0, T]$$

$$z_i(\tau) \in Z_i, u_i(\tau) \in U_i, \quad \forall i, \tau \in [0, T]$$

$$\sum_{i=1}^N H_i z_i(\tau) \leq h, \quad \tau \in [0, T]$$

$$\sum_{i=1}^N g_i(x_i) \leq 0$$

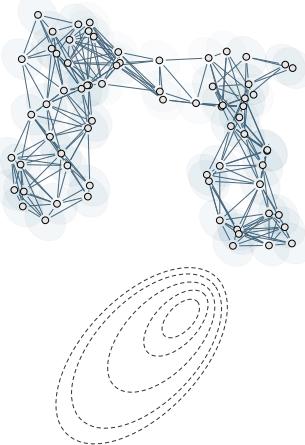


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Challenge Wrap-up

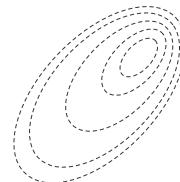
Network

- Undirected/Directed
- Fixed/Time-Varying
- Asynchronous, Unreliable



Optimization

- Nonsmooth (convex), Nonconvex, Mixed-Integer, Combinatorial
- Big-Data (high dim. dec. var.)
- Dynamic, Online, Stochastic, Uncertain



Selected Topics

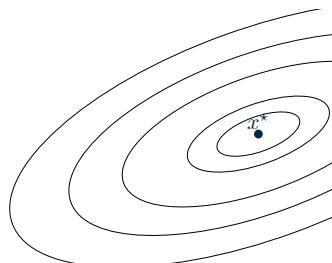
Cost coupled - Distributed Big-Data Optimization

Common cost - Constraint Exchange

Constraint coupled - Distributed Primal Decomposition

Consensus-based Method

$$\min_x \sum_{i=1}^N f_i(x)$$



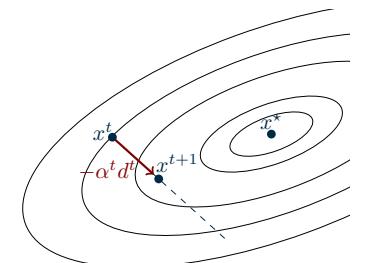
Consensus-based Method

$$\min_x \sum_{i=1}^N f_i(x)$$

Centralized (sub)gradient

$$x^{t+1} = x^t - \alpha^t d$$

d
update direction



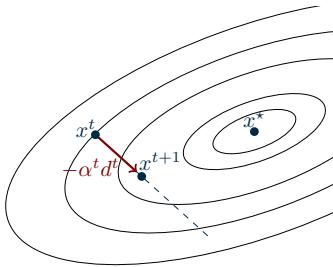
Consensus-based Method

$$\min_x \sum_{i=1}^N f_i(x)$$

Centralized (sub)gradient

$$x^{t+1} = x^t - \alpha^t \left[\sum_{h=1}^N \nabla f_h(x^t) \right]$$

update direction



Consensus-based Method

$$\min_x \sum_{i=1}^N f_i(x)$$

Centralized (sub)gradient

$$x^{t+1} = x^t - \alpha^t \left[\sum_{h=1}^N \nabla f_h(x^t) \right]$$

update direction



Distributed update

$$x_i^{t+1} =$$

Consensus-based Method

$$\min_x \sum_{i=1}^N f_i(x)$$

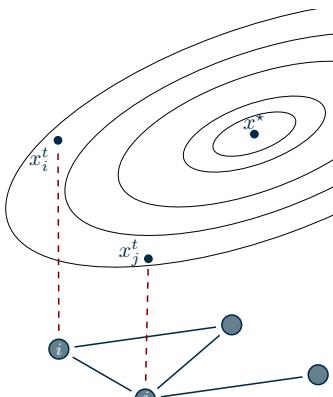
Centralized (sub)gradient

$$x^{t+1} = x^t - \alpha^t \left[\sum_{h=1}^N \nabla f_h(x^t) \right]$$

update direction

Distributed update

$$x_i^{t+1} =$$



Consensus-based Method

$$\min_x \sum_{i=1}^N f_i(x)$$

Centralized (sub)gradient

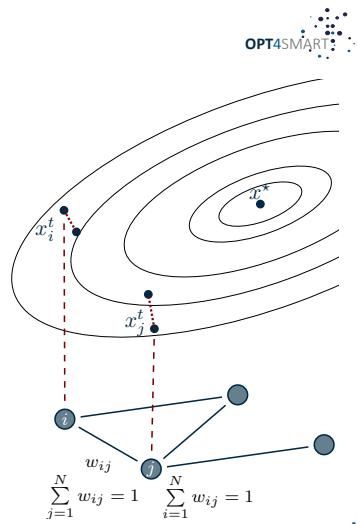
$$x^{t+1} = x^t - \alpha^t \left[\sum_{h=1}^N \nabla f_h(x^t) \right]$$

update direction

Distributed update

$$x_i^{t+1} = \left[\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t \right]$$

averaging



Consensus-based Method

$$\min_x \sum_{i=1}^N f_i(x)$$

Centralized (sub)gradient

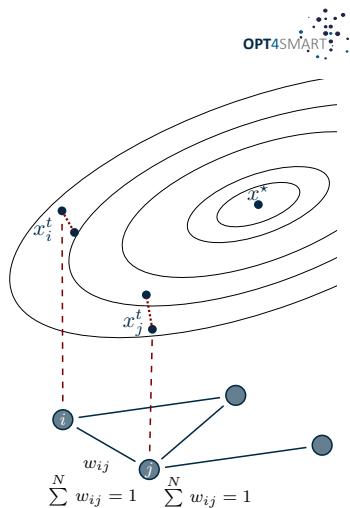
$$x^{t+1} = x^t - \alpha^t \sum_{h=1}^N \nabla f_h(x^t)$$

update direction

Distributed update

$$x_i^{t+1} = \sum_{j \in \mathcal{N}_i} w_{ij} x_j^t$$

averaging
(push-sum for digraphs)



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Consensus-based Method

$$\min_x \sum_{i=1}^N f_i(x)$$

Centralized (sub)gradient

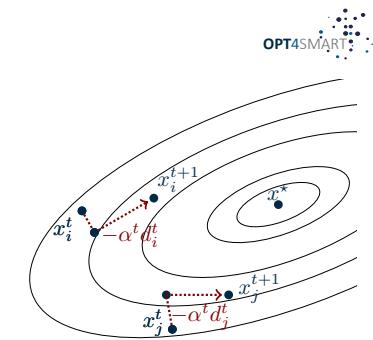
$$x^{t+1} = x^t - \alpha^t \sum_{h=1}^N \nabla f_h(x^t)$$

global

Distributed update

$$x_i^{t+1} = \sum_{j \in \mathcal{N}_i} w_{ij} x_j^t - \alpha^t d_i^t$$

averaging
local



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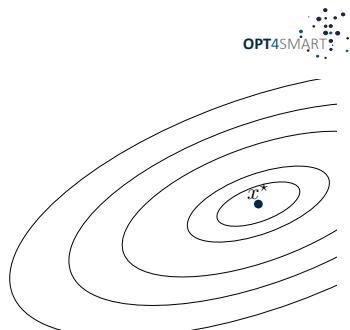
Distributed (Sub)gradient

$$\min_x \sum_{i=1}^N f_i(x)$$

Distributed (sub)gradient

$$x_i^{t+1} = \sum_{j \in \mathcal{N}_i} w_{ij} x_j^t - \alpha^t \nabla f_i \left(\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t \right)$$

d_i^t



Nedić & Ozdaglar, "Distributed subgradient methods for multi-agent optimization." TAC 2009

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Distributed (Sub)gradient

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Distributed (sub)gradient

$$x_i^{t+1} = \sum_{j \in \mathcal{N}_i} w_{ij} x_j^t - \alpha^t \nabla f_i \left(\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t \right)$$

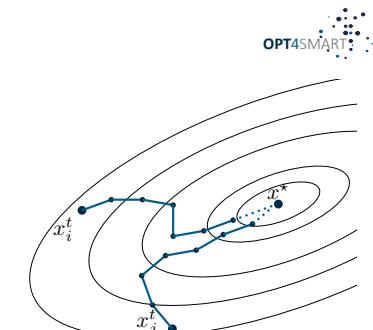
Theorem

Convex costs f_i, \dots

Diminishing stepsize ($\alpha^t \rightarrow 0$ and ...)

Doubly-stochastic weights

\Rightarrow Consensus: $\lim_{t \rightarrow \infty} \|x_i^t - \bar{x}^t\| = 0$
Optimality: $\lim_{t \rightarrow \infty} \|\bar{x}^t - x^*\| = 0$



Nedić & Ozdaglar, "Distributed subgradient methods for multi-agent optimization." TAC 2009

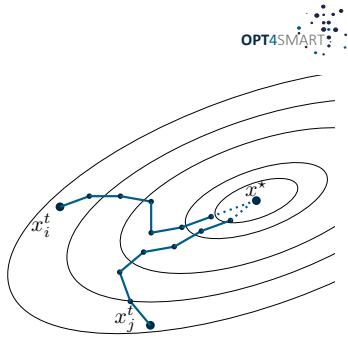
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Distributed (Sub)gradient

$$\min_x \sum_{i=1}^N f_i(x)$$

Distributed (sub)gradient

$$x_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t} - \alpha^t \boxed{\nabla f_i \left(\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t \right)}$$



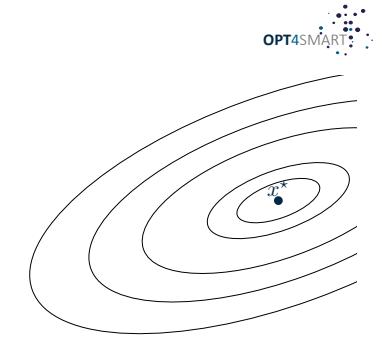
Some early literature

[Johansson,CDC'08], [Nedic,TAC'09], [Nedic,TAC'10], [Cattivelli,TSP'10],
 [Ram,JOTA'10], [Lobel,TAC'11], [Wang,CDC'11], [Chen,TSP'12], [Lu,TAC'12],
 [Bianchi,TAC'13], [Lee,STSP'13], [Jakovetic,TAC'14], [Shi,SJO'15],
 [Shi,TSP'15], [Nedic,TAC'15], ...

Gradient Tracking

$$\min_x \sum_{i=1}^N f_i(x)$$

$$x_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t} - \alpha \boxed{d_i^t}$$



Gradient Tracking

$$\min_x \sum_{i=1}^N f_i(x)$$

$$x_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t} - \alpha \boxed{d_i^t}$$



Main idea:

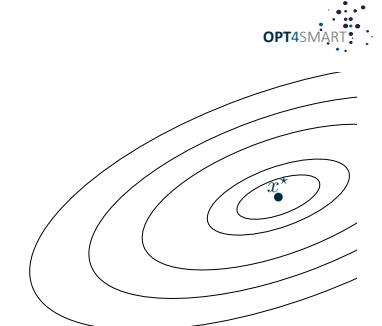
$$d_i^t \xrightarrow{t \rightarrow \infty} \boxed{\frac{1}{N} \sum_{h=1}^N \nabla f_h(x_h^t)}$$

average of gradients

Gradient Tracking

$$\min_x \sum_{i=1}^N f_i(x)$$

$$x_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t} - \alpha \boxed{d_i^t}$$



Dynamic average consensus

$$d_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} \tilde{w}_{ij} d_j^t} + \boxed{(\nabla f_i(x_i^{t+1}) - \nabla f_i(x_i^t))}$$

Gradient Tracking

$$\min_x \sum_{i=1}^N f_i(x)$$

$$x_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t} - \alpha \boxed{d_i^t}$$

$$d_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} \tilde{w}_{ij} d_j^t} + \boxed{(\nabla f_i(x_i^{t+1}) - \nabla f_i(x_i^t))}$$



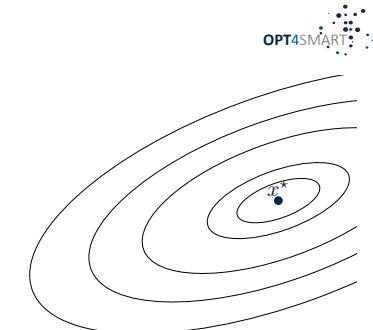
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Gradient Tracking

$$\min_x \sum_{i=1}^N f_i(x)$$

$$x_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t} - \alpha \boxed{d_i^t}$$

$$d_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} \tilde{w}_{ij} d_j^t} + \boxed{(\nabla f_i(x_i^{t+1}) - \nabla f_i(x_i^t))}$$



Theorem

Smooth (nonconvex) costs f_i, \dots
Constant stepsize α
Row (w_{ij}) Clmn (\tilde{w}_{ij}) stoch. weights

Consensus: $\lim_{t \rightarrow \infty} \|x_i^t - \bar{x}^t\| = 0$
Optimality: $\lim_{t \rightarrow \infty} \|\bar{x}^t - x^*\| = 0$

linear rate

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Gradient Tracking

$$\min_x \sum_{i=1}^N f_i(x)$$

$$x_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} w_{ij} x_j^t} - \alpha \boxed{d_i^t}$$

$$d_i^{t+1} = \boxed{\sum_{j \in \mathcal{N}_i} \tilde{w}_{ij} d_j^t} + \boxed{(\nabla f_i(x_i^{t+1}) - \nabla f_i(x_i^t))}$$



Some history of gradient tracking

[Zanella,CDC-ECC'11], [DiLorenzo,CAMSAP'15], [Xu,CDC'15],
[Nedic,CDC'16], [Qu,CDC'16], [Sun,Asilomar'16], [Qu,CDC'17], [Xi,TAC'18],
[Xu,TAC'18], [Xin,L-CSS'18], ...

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Other Approaches for Cost Coupled

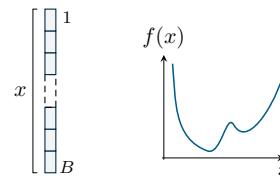
- Dual decomposition
- Primal-Dual
- Alternating Direction Method of Multipliers (ADMM)



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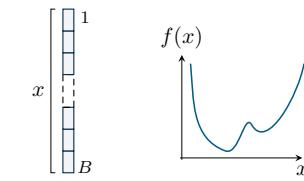
Block-wise Gradient Tracking

$$\min_x \sum_{i=1}^N f_i(x)$$



Block-wise Gradient Tracking

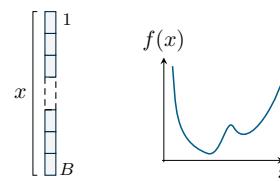
$$\min_x \sum_{i=1}^N f_i(x)$$



collaboration with G. Scutari and Y. Sun (Purdue University)

Block-wise Gradient Tracking

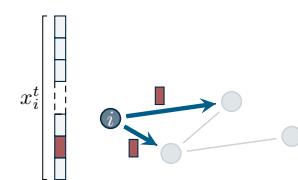
$$\min_x \sum_{i=1}^N f_i(x)$$



minimizing w.r.t. x too costly
transmitting x unaffordable

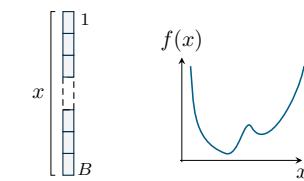


block-wise
optimization/communication



Block-wise Gradient Tracking

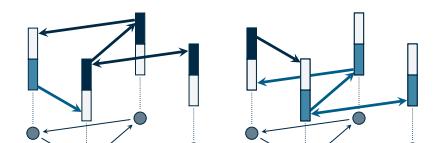
$$\min_x \sum_{i=1}^N f_i(x)$$



Block selection rule:
 ℓ_i^t block selected at time t by agent i

Block-dependent neighbor set:
neighbors of i sending block ℓ at time t

$$\mathcal{N}_{i,\ell}^t \triangleq \{j \in \mathcal{N}_i \mid \ell_j^t = \ell\} \cup \{i\} \subseteq \mathcal{N}_i$$



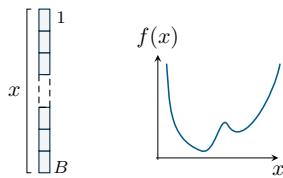
Block-wise Gradient Tracking

$$\min_x \sum_{i=1}^N f_i(x)$$

$$\Delta x_{(i,\ell)}^t = \begin{cases} -d_{(i,\ell)}^t, & \text{if } \ell = \ell_i^t \\ 0, & \text{otherwise} \end{cases}$$

$$x_{(i,\ell)}^{t+1} = \sum_{j \in N_{i,\ell}^t} \frac{w_{ij\ell}^t \phi_{(j,\ell)}^t}{\phi_{(i,\ell)}^{t+1}} x_{(j,\ell)}^t + \alpha^t \phi_{(i,\ell)}^t \Delta x_{(i,\ell)}^t$$

$$d_{(i,\ell)}^{t+1} = \sum_{j \in N_{i,\ell}^t} \frac{w_{ij\ell}^t \phi_{(j,\ell)}^t}{\phi_{(i,\ell)}^{t+1}} d_{(j,\ell)}^t + \frac{\nabla_\ell f_i(x_{(i,:)}^{t+1}) - \nabla_\ell f_i(x_{(i,:)}^t)}{\phi_{(i,\ell)}^{t+1}}$$



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Block-wise Gradient Tracking Convergence

Theorem

Let $\{x_1^t, \dots, x_N^t\}_{t \geq 0}$ be generated by block-wise gradient tracking with

- strongly connected graph + ess. cyclic block-selection rule
- diminishing step-size α^t
- smooth (nonconvex) costs f_i, \dots

Let $\bar{x}^t \triangleq \frac{1}{N} \sum_{i=1}^N \phi_{(i,:)}^t x_{(i,:)}^t$, then

- (i) consensus: $\|x_{(i,:)}^t - \bar{x}^t\| \rightarrow 0$ as $t \rightarrow \infty$, for all $i \in \{1, \dots, N\}$;
- (ii) convergence: every limit point $\{\bar{x}^t\}_{t \geq 0}$ is a stationary solution.

Notarnicola et al., “Distributed Big-Data Optimization via Block-wise Gradient Tracking.” arXiv 1808.07252

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Block-wise Gradient Tracking Convergence

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Extensions: regularization and constraints, single-block gradient update

Notarnicola et al., “Distributed Big-Data Optimization via Block-wise Gradient Tracking.” arXiv 1808.07252

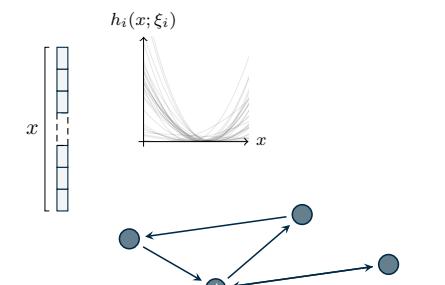
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Stochastic Big-Data Optimization

$$\min_x \sum_{i=1}^N \mathbb{E}_{\xi_i} [h_i(x; \xi_i)]$$

Problem features:

- Big-data
- Stochastic
- Nonsmooth (Convex)



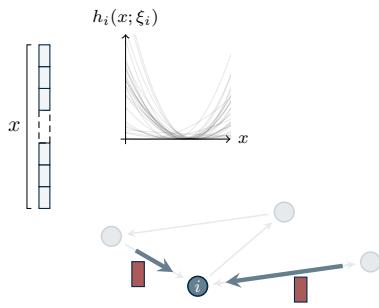
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Stochastic Big-Data Optimization

$$\min_x \sum_{i=1}^N \mathbb{E}_{\xi_i} [h_i(x; \xi_i)]$$

Distributed Block Subgradient

$$y_i^t = \sum_{j \in \mathcal{N}_i} w_{ij} x_j^t$$



Stochastic Big-Data Optimization

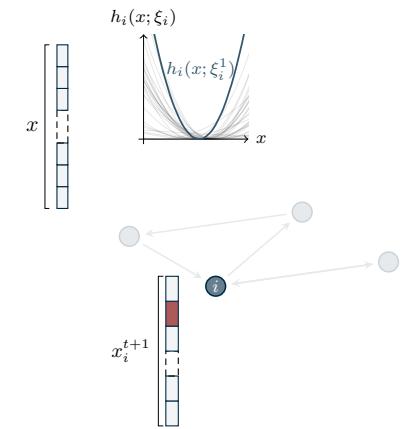
$$\min_x \sum_{i=1}^N \mathbb{E}_{\xi_i} [h_i(x; \xi_i)]$$

Distributed Block Subgradient

$$y_i^t = \sum_{j \in \mathcal{N}_i} w_{ij} x_j^t$$

DRAW ℓ_i^t and ξ_i^t

$$x_{i,\ell}^{t+1} = \begin{cases} y_{i,\ell}^t - \alpha_i^t [\nabla h_i(y_i^t; \xi_i^t)]_\ell, & \text{if } \ell = \ell_i^t \\ x_{i,\ell}^t, & \text{otherwise} \end{cases}$$



Stochastic Big-Data Optimization

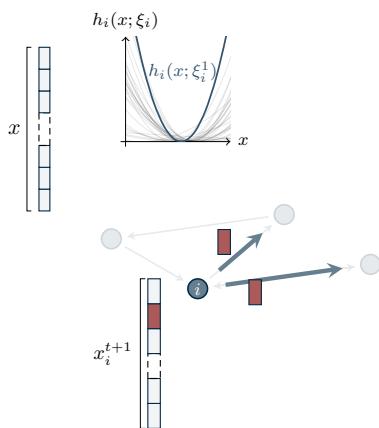
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Distributed Block Subgradient

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DRAW ℓ_i^t and ξ_i^t

$$x_{i,\ell}^{t+1} = \begin{cases} y_{i,\ell}^t - \alpha_i^t [\nabla h_i(y_i^t; \xi_i^t)]_\ell, & \text{if } \ell = \ell_i^t \\ x_{i,\ell}^t, & \text{otherwise} \end{cases}$$



Distributed Block Subgradient Convergence

Theorem

Let $\{x_1^t, \dots, x_N^t\}_{t \geq 0}$ be generated by distributed block subgradient with

- α_i^t diminishing
- doubly stochastic weights
- unbiased stochastic subgradients

Then, for all $i \in \{1, \dots, N\}$

$$f_{\text{best}}(x_i^t) \xrightarrow[t \rightarrow \infty]{} f^*$$

Distributed Block Subgradient Convergence

Theorem

Let $\{x_1^t, \dots, x_N^t\}_{t \geq 0}$ be generated by distributed block subgradient with

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Then, for all $i \in \{1, \dots, N\}$

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Extensions

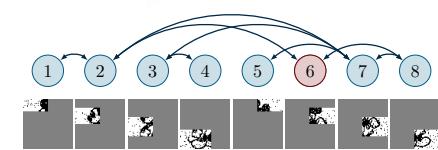
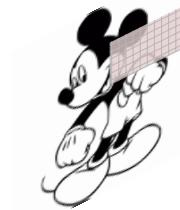
- constraint set $X \subseteq \mathbb{R}^d$ and proximal mapping
- awake/idle nodes

Farina & Notarstefano, "Randomized Block Proximal Methods for Distributed Stochastic Big-Data Optimization" arXiv 2019

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Distributed Submodular Minimization

$$\min_{X \subseteq V} \sum_{i=1}^N F_i(X)$$



Farina, Testa & Notarstefano, "Distributed Submodular Minimization via Block-Wise Updates and Communications" arXiv 2019

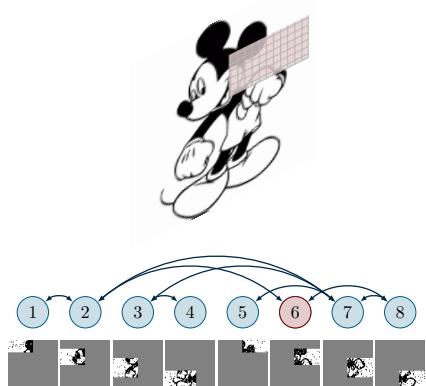
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Distributed Submodular Minimization

$$\min_{X \subseteq V} \sum_{i=1}^N F_i(X)$$

↓
Lovàsz extension

$$\min_{x \in [0,1]^{|V|}} \sum_{i=1}^N f_i(x)$$



Farina, Testa & Notarstefano, "Distributed Submodular Minimization via Block-Wise Updates and Communications" arXiv 2019

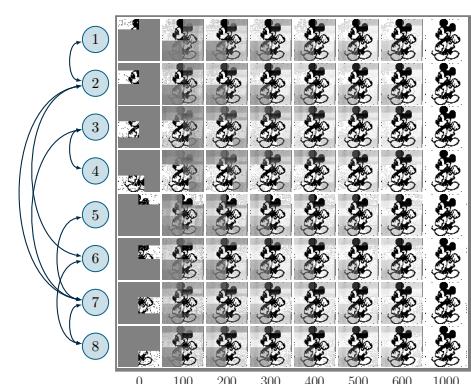
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Distributed Submodular Minimization

$$\min_{X \subseteq V} \sum_{i=1}^N F_i(X)$$

↓
Lovàsz extension

$$\min_{x \in [0,1]^{|V|}} \sum_{i=1}^N f_i(x)$$



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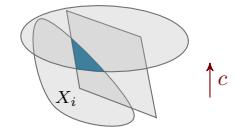
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Cost coupled - Distributed Big-Data Optimization

Common cost - Constraint Exchange

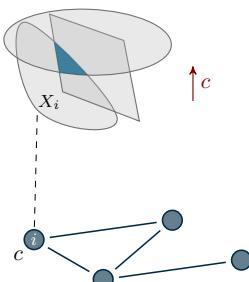
Constraint coupled - Distributed Primal Decomposition

$$\begin{aligned} & \min_x c^\top x \\ \text{subj.to } & x \in \bigcap_{i=1}^N X_i \end{aligned}$$

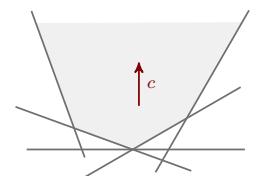


$$\begin{aligned} & \min_x c^\top x \\ \text{subj.to } & x \in \bigcap_{i=1}^N X_i \end{aligned}$$

- agent i knows c and X_i
- convex optimization problem
(abstract programs)
- asynchronous, unreliable, directed communication

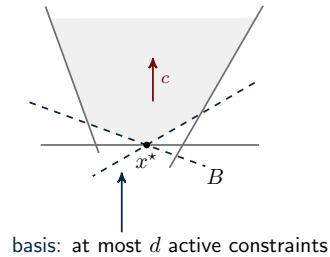


$$\begin{aligned} & \min_x c^\top x \\ \text{subj.to } & \underbrace{a_i^\top x \leq b_i}_{X_i} \quad i \in \{1, \dots, N\} \end{aligned}$$



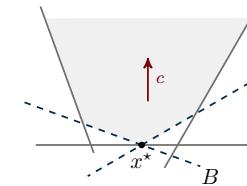
Intuition on Linear Programming

$$\begin{aligned} & \min_x c^T x \\ \text{subj.to } & \underbrace{a_i^T x \leq b_i}_{X_i} \quad i \in \{1, \dots, N\} \end{aligned}$$



Intuition on Linear Programming

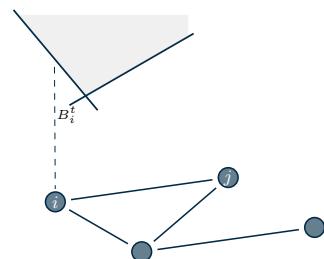
$$\begin{aligned} & \min_x c^T x \\ \text{subj.to } & \underbrace{a_i^T x \leq b_i}_{X_i} \quad i \in \{1, \dots, N\} \end{aligned}$$



Key idea: each agent stores and exchanges a candidate solution basis B_i^t

Constraints Consensus Algorithm

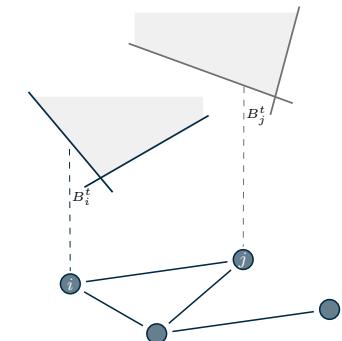
Initial constraint: X_i



Constraints Consensus Algorithm

Initial constraint: X_i

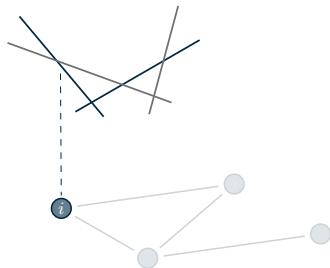
Gather B_j^t from $j \in \mathcal{N}_i^t$



Constraints Consensus Algorithm

Initial constraint: X_i

Gather B_j^t from $j \in \mathcal{N}_i^t$



Constraints Consensus Algorithm

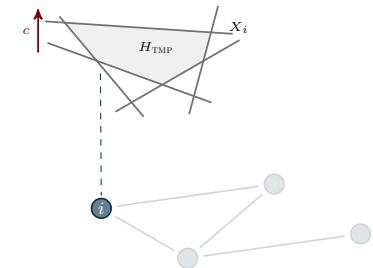
Initial constraint: X_i

Gather B_j^t from $j \in \mathcal{N}_i^t$

Build H_{TMP} and compute x_i^{t+1} as

$$\underset{x}{\operatorname{argmin}} \quad c^\top x$$

subj.to $x \in H_{\text{TMP}}$



Constraints Consensus Algorithm

Initial constraint: X_i

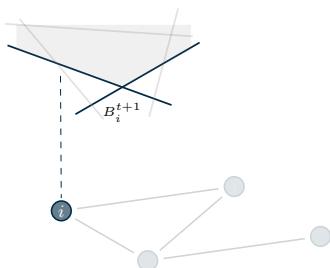
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Update B_i^{t+1} as a basis of x_i^{t+1}



Constraints Consensus Algorithm

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Theorem

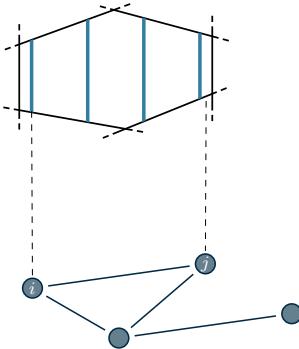
Jointly strongly connected digraph

Abstract programs (e.g., convex programs)

⇒ There exists $T > 0$ s.t. $x_i^t = x^*, \forall i, \forall t > T$

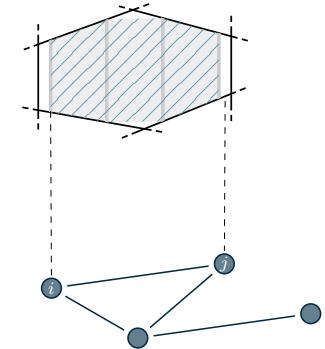
Mixed Integer Linear Programming

$$\begin{aligned} & \min_x c^\top x \\ \text{subj.to } & a_i^\top x \leq b_i \quad i \in \{1, \dots, N\} \\ & x \in \mathbb{Z}^{d_Z} \times \mathbb{R}^{d_R} \end{aligned}$$



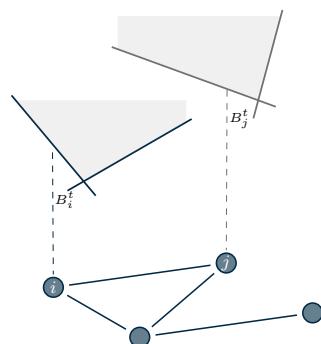
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Cut Generation and Constraint Exchange

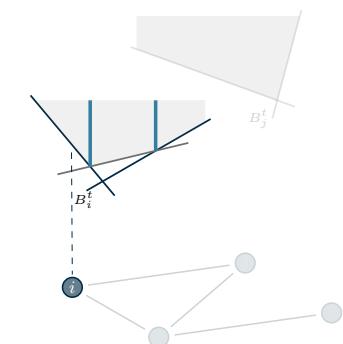
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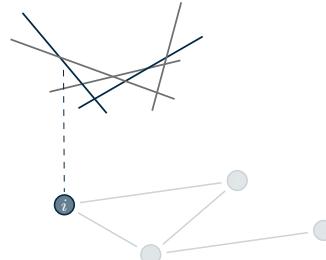
Cut Generation and Constraint Exchange

Initial constraint: X_i

Generate cutting planes
(Gomory and cost-based cuts)



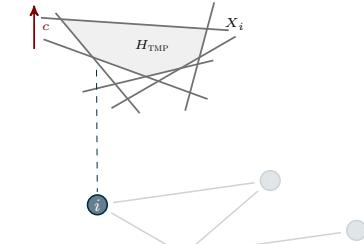
Cut Generation and Constraint Exchange

Initial constraint: X_i Generate cutting planes
(Gomory and cost-based cuts)Gather B_j^t from $j \in \mathcal{N}_i^t$ 

Cut Generation and Constraint Exchange

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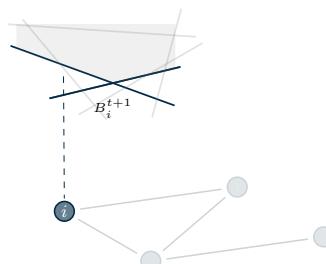
Build H_{TMP} and compute x_i^{t+1} as

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Cut Generation and Constraint Exchange

Initial constraint: X_i Generate cutting planes
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Update B_i^{t+1} as a basis of x_i^{t+1} 

Distributed Algorithm for Integer-Valued MILPs

Theorem

Let $\{x_1^t, \dots, x_N^t\}_{t \geq 0}$ be generated by cut generation and constraint exchange algorithm with

- jointly strongly connected digraph
- bounded feasible set
- integer optimal cost

Then, there exists $T > 0$ such that $\forall i = 1, \dots, N$

$$x_t^i = x^*, \quad \text{for all } t \geq T \quad (\text{optimal solution})$$

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Extensions

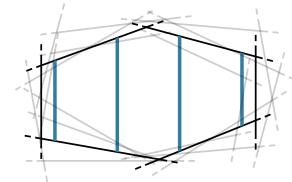
- general MILPs
- ϵ -suboptimal solution by approximate epigraph reformulation

Testa et al., "Distributed Mixed-Integer Linear Programming via Cut Generation and Constraint Exchange." TAC 2020 (in press)

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Distributed Robust Optimization

$$\begin{aligned} & \min_x c^\top x \\ \text{subj.to } & x \in \bigcap_{i=1}^N X_i(q), \forall q \in \mathbb{Q} \\ & x \in \mathbb{Z}^{d_Z} \times \mathbb{R}^{d_R} \end{aligned}$$



collaboration with
M. Chamanbaz and R. Bouffanais (SUTD)

- convex (mixed-integer) robust programs
- agents know cost c and $X_i(q)$
($q \in \mathbb{Q}$ uncertain parameters)
- local verification and re-optimization

Chamanbaz et al., "Randomized Constraints Consensus for Distributed Robust Mixed-Integer Programming." (submitted)

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Selected Topics

Cost coupled - Distributed Big-Data Optimization

Common cost - Constraint Exchange

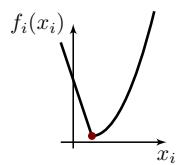
Constraint coupled - Distributed Primal Decomposition

Constraint-Coupled Optimization

$$\begin{aligned} & \min_{x_1, \dots, x_N} \sum_{i=1}^N f_i(x_i) \\ \text{subj.to } & x_i \in X_i \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b \end{aligned}$$

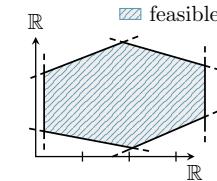
Constraint-Coupled Optimization

$$\begin{aligned} \min_{x_1, \dots, x_N} \quad & \sum_{i=1}^N f_i(x_i) && \xleftarrow{\text{convex functions}} \\ \text{subj.to} \quad & x_i \in X_i \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b \end{aligned}$$



Constraint-Coupled Optimization

$$\begin{aligned} \min_{x_1, \dots, x_N} \quad & \sum_{i=1}^N f_i(x_i) \\ \text{subj.to} \quad & x_i \in X_i \quad \forall i && \xleftarrow{\text{convex, compact local sets}} \\ & \sum_{i=1}^N A_i x_i \leq b \end{aligned}$$



Constraint-Coupled Optimization

$$\begin{aligned} \min_{x_1, \dots, x_N} \quad & \sum_{i=1}^N f_i(x_i) \\ \text{subj.to} \quad & x_i \in X_i \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b && \xleftarrow{\text{linear coupling constraints}} \end{aligned}$$

b

$A_1 x_1$	$A_i x_i$	$A_N x_N$
-----------	-------	-----------	-------	-----------

Constraint-Coupled Optimization

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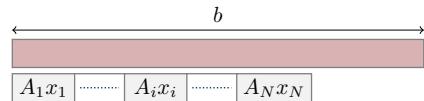
Some (incomplete) literature

Duality-based: [Bürger,TAC'14], [Chang,TAC'14], [Simonetto,JOTA'16], [Falsone,Aut'17], ...

Resource allocation: [Lakshmanan,SJO'08], [Necoara,TAC'13], [Cherukuri,TCNS'15], ...

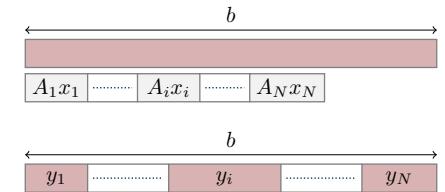
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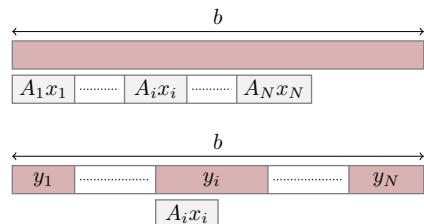
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Constraint-Coupled Optimization

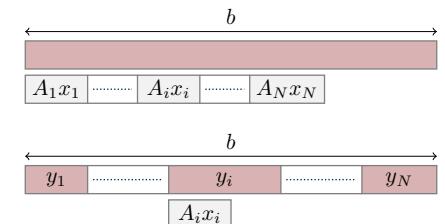
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$$\begin{aligned} & \min_{x_i} f_i(x_i) \\ \text{subj.to } & x_i \in X_i \\ & A_i x_i \leq y_i \end{aligned}$$

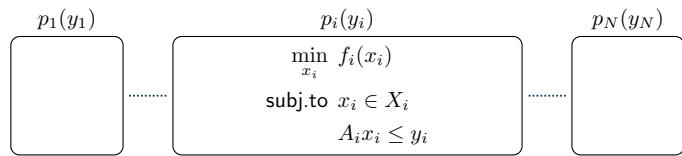
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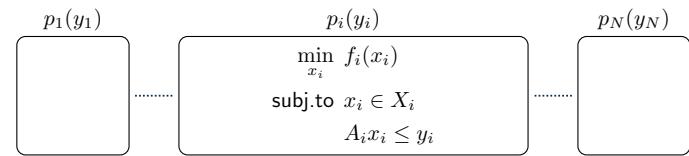
$$p_i(y_i) = \boxed{\begin{aligned} & \min_{x_i} f_i(x_i) \\ \text{subj.to } & x_i \in X_i \\ & A_i x_i \leq y_i \end{aligned}}$$

Primal Decomposition



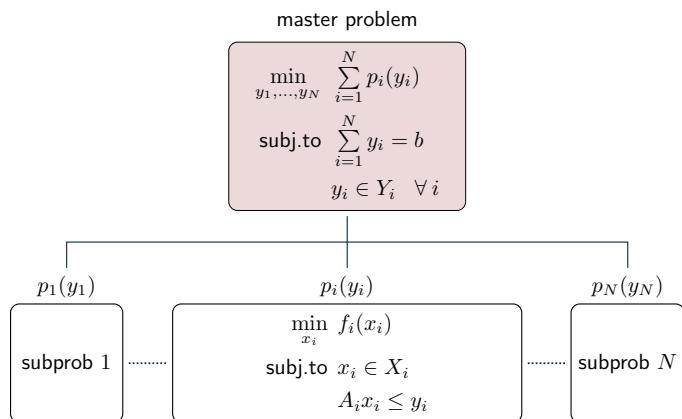
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$$\begin{aligned} & \min_{y_1, \dots, y_N} \sum_{i=1}^N p_i(y_i) \\ \text{subj.to } & \sum_{i=1}^N y_i = b \\ & y_i \in Y_i \quad \forall i \end{aligned}$$



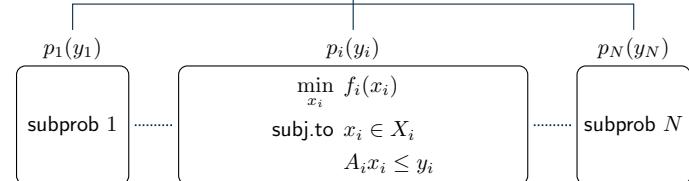
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Primal Decomposition



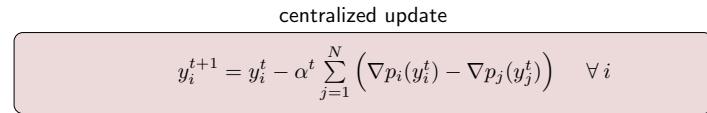
Giuseppe Notarstefano – Optimization in Networkland: Challenges and Opportunities – ECC '19, Naples – 33

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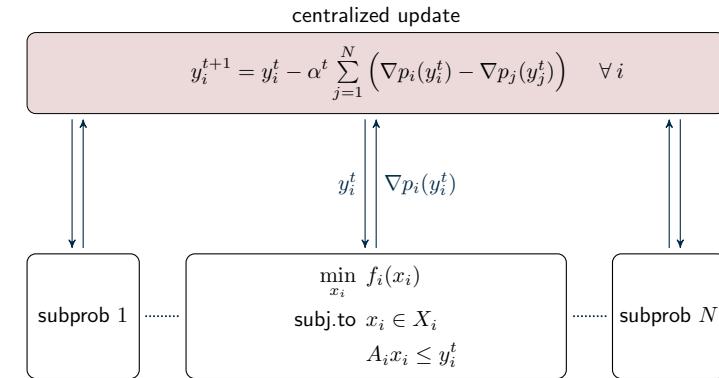


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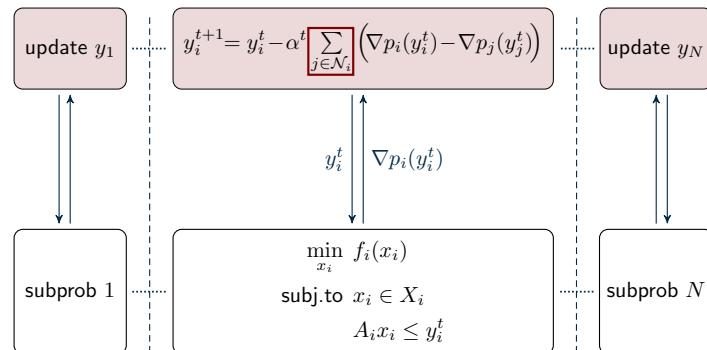
Subgradient Method on Master Problem



Subgradient Method on Master Problem



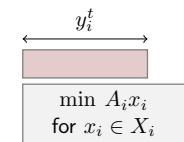
Subgradient Method on Master Problem



More on Subproblems

Q1: What if subproblem is infeasible?

$$\begin{aligned} & \min_{x_i} f_i(x_i) \\ \text{subj.to } & x_i \in X_i \\ & A_i x_i \leq y_i^t \end{aligned}$$



infeasible local problem

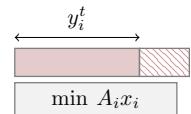
Q2: How to compute $\nabla p_i(y_i^t)$?

More on Subproblems

More on Subproblems

Q1: What if subproblem is infeasible?

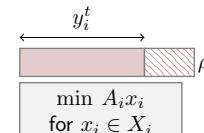
$$\begin{array}{ll} \min_{x_i, \rho_i} & f_i(x_i) + M\rho_i \\ \text{subj.to} & x_i \in X_i, \quad \rho_i \geq 0 \\ & A_i x_i \leq y_i^t + \rho_i \mathbf{1} \end{array}$$



transient violation

Q1: What if subproblem is infeasible?

$$\begin{array}{ll} \min_{x_i, \rho_i} & f_i(x_i) + M\rho_i \\ \text{subj.to} & x_i \in X_i, \quad \rho_i \geq 0 \\ & \mu_i : A_i x_i \leq y_i^t + \rho_i \mathbf{1} \end{array}$$



transient violation

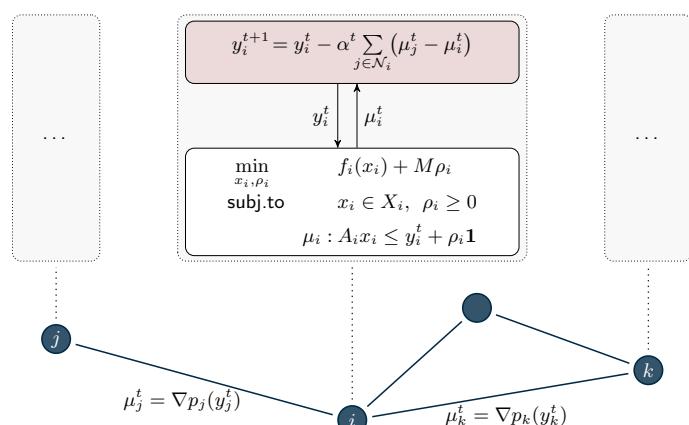
Q2: How to compute $\nabla p_i(y_i^t)$?

Q2: How to compute $\nabla p_i(y_i^t)$?

Use multiplier μ_i^t of $A_i x_i \leq y_i^t + \rho_i \mathbf{1}$:

$$\nabla p_i(y_i^t) = -\mu_i^t$$

Distributed Primal Decomposition Algorithm



Distributed Primal Decomposition Convergence

Theorem

Let $\{x_1^t, \dots, x_N^t\}_{t \geq 0}$ be generated by distributed algorithm with

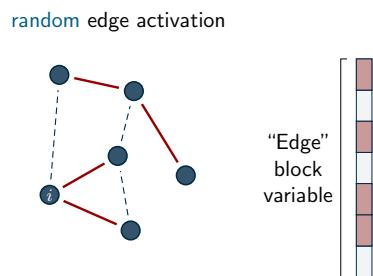
- α^t diminishing step-size
- $M > 0$ sufficiently large

Then:

- $\lim_{t \rightarrow \infty} \sum_{i=1}^N f_i(x_i^t) = f^*$
- $\{x_1^t, \dots, x_N^t\}_{t \geq 0} \xrightarrow{\text{every limit point}} (x_1^*, \dots, x_N^*)$ optimal solution

Extension to Time-varying Graphs

$$y_i^{t+1} = y_i^t - \alpha^t \sum_{j \in \mathcal{N}_i^t} (\mu_j^{t+1} - \mu_i^{t+1})$$



Analysis approach

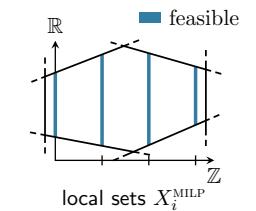
- Block subgradient method
- Random block (edge) selection

Camisa et al., "Distributed Constraint-Coupled Optimization over Random Time-Varying Graphs via Primal Decomposition and Block Subgradient Approaches." (subm. to conf.)

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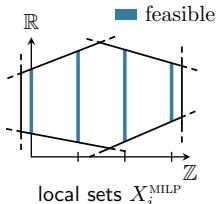
Mixed-Integer Linear Programs

$$\begin{aligned} & \min_{x_1, \dots, x_N} \sum_{i=1}^N c_i^\top x_i \\ \text{subj.to } & x_i \in X_i^{\text{MILP}} \subset \mathbb{Z}^{z_i} \times \mathbb{R}^{r_i} \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b \end{aligned}$$



Mixed-Integer Linear Programs

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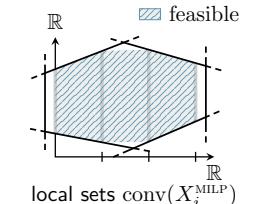
Challenge: large-scale and NP-hard problem to be solved in short time

Goal: fast computation of "high-quality" suboptimal solutions

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Mixed-Integer Linear Programs

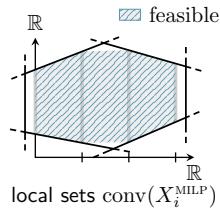
$$\begin{aligned} & \min_{x_1, \dots, x_N} \sum_{i=1}^N c_i^\top x_i \\ \text{subj.to } & x_i \in \text{conv}(X_i^{\text{MILP}}) \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b \end{aligned}$$



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Mixed-Integer Linear Programs

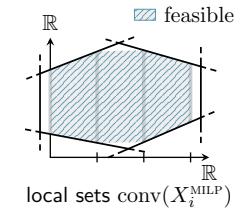
$$\begin{aligned} \min_{x_1, \dots, x_N} \quad & \sum_{i=1}^N c_i^\top x_i \\ \text{subj.to} \quad & x_i \in \text{conv}(X_i^{\text{MILP}}) \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b \quad \leftarrow S \text{ constraints} \end{aligned}$$



Theorem (Shapley-Folkman): Let $(x_1^{\text{conv}}, \dots, x_N^{\text{conv}})$ be unique optimal solution
Then x_i^{conv} already mixed integer for at least $N - S$ agents

Mixed-Integer Linear Programs

$$\begin{aligned} \min_{x_1, \dots, x_N} \quad & \sum_{i=1}^N c_i^\top x_i \\ \text{subj.to} \quad & x_i \in \text{conv}(X_i^{\text{MILP}}) \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b \quad \leftarrow S \text{ constraints} \end{aligned}$$

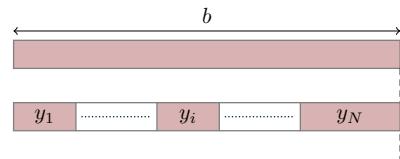


Theorem (Shapley-Folkman): Let $(x_1^{\text{conv}}, \dots, x_N^{\text{conv}})$ be unique optimal solution
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State of the art: dual decomposition (Vujanic,Aut'16, Falsone,Aut'19)

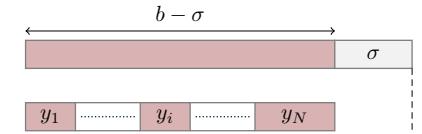
Restriction of Coupling Constraints

$$\begin{aligned} \min_{x_1, \dots, x_N} \quad & \sum_{i=1}^N c_i^\top x_i \\ \text{subj.to} \quad & x_i \in \text{conv}(X_i^{\text{MILP}}) \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b \end{aligned}$$



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Restriction of Coupling Constraints

$b - \sigma$

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Distributed Primal Decomposition for MILP



Convexified problem

$$\begin{aligned} & \min_{x_1, \dots, x_N} \sum_{i=1}^N c_i^\top x_i \\ \text{subj.to } & x_i \in \text{conv}(X_i^{\text{MILP}}) \quad \forall i \\ & \sum_{i=1}^N A_i x_i \leq b - \sigma \end{aligned}$$

Solve with Distributed Primal Decomposition

$$y_i^t \rightarrow y_i^{\text{conv}}$$

Distributed Primal Decomposition for MILP



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Solve with Distributed Primal Decomposition

$$y_i^t \rightarrow y_i^{\text{conv}}$$

$$\begin{aligned} & \min A_i x_i \\ \text{for } & x_i \in X_i^{\text{MILP}} \end{aligned}$$

Compute $x_i^t \in X_i^{\text{MILP}}$ with minimal violation of $A_i x_i \leq y_i^t$

Finite-time Theoretical Results

Theorem

Let $\{x_1^t, \dots, x_N^t\}_{t \geq 0}$ be generated by distributed algorithm for MILP.

Finite-time Theoretical Results

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Let $\{x_1^t, \dots, x_N^t\}_{t \geq 0}$ be generated by distributed algorithm for MILP.

Then, for given "extra restriction", $\exists T > 0$ such that

- MILP feasibility:

$$x_i^t \in X_i^{\text{MILP}} \quad \forall i, \quad \sum_{i=1}^N A_i x_i^t \leq b, \quad \forall t \geq T$$

Camisa & Notaricola & Notarstefano, "A Primal Decomposition Method with Suboptimality Bounds for Distributed Mixed-Integer Linear Programming." CDC 2018

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- MILP suboptimality bound:

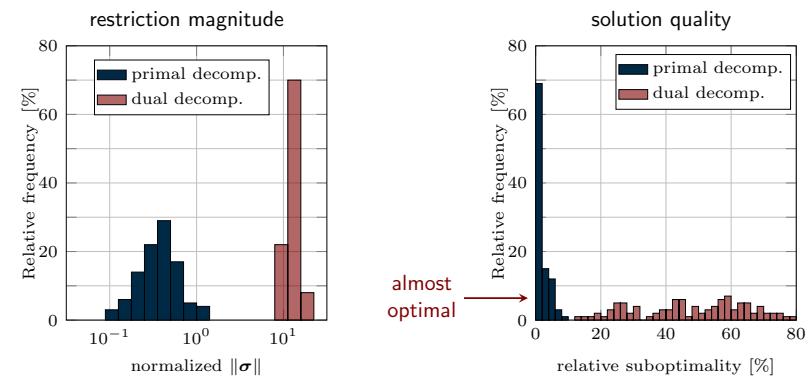
$$\sum_{i=1}^N c_i^\top x_i^t - J^{\text{MILP}} \leq \begin{array}{c} \text{convexification} \\ \text{suboptimality} \end{array} + \begin{array}{c} \text{restriction} \\ \text{suboptimality} \end{array} + \begin{array}{c} \text{distance to} \\ \text{convergence} \end{array} \quad \forall t \geq T$$

Camisa & Notaricola & Notarstefano, "A Primal Decomposition Method with Suboptimality Bounds for Distributed Mixed-Integer Linear Programming." CDC 2018

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Numerical Computations

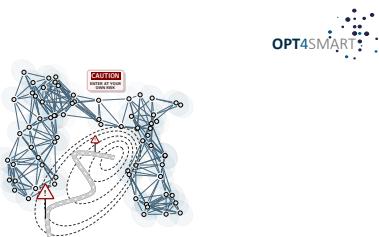
Montecarlo simulations with $X_i \subset \mathbb{Z} \times \mathbb{R}$ (100 instances, $N = 50$, $S = 10$)



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What Next?

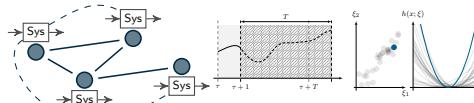
Network/optimization challenges



System theoretical approach
to distributed optimization

$$\begin{array}{c} u \\ \downarrow \\ \left[\begin{array}{c} x_1^{k+1} \\ x_2^{k+1} \\ \vdots \\ x_N^{k+1} \end{array} \right] = W \left[\begin{array}{c} x_1^k \\ x_2^k \\ \vdots \\ x_N^k \end{array} \right] - \gamma u \\ \downarrow \\ \left[\begin{array}{c} x_1^{k+1} \\ x_2^{k+1} \\ \vdots \\ x_N^{k+1} \end{array} \right] = \left[\begin{array}{c} x_1^k \\ x_2^k \\ \vdots \\ x_N^k \end{array} \right] + (\bar{W} - I) \left[\begin{array}{c} \nabla f_1(x_1) \\ \vdots \\ \nabla f_N(x_N) \end{array} \right] \end{array}$$

Optimal control of complex systems
model-based vs data-driven



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PEOPLE PUBLICATIONS GUESTBOOK

NEWS

29 MAY 2018 Bologna
Brought to you by Dr. Andrea Sironi (IBM research, USA) and Prof. Alessandro Testa (Personalized Optimal Control in a Time-varying World)

20 MAY 2018 Bologna
Our former Post-doc Alessandro Fazio is actively involved in the development of a system for driverless driving with Vihab. A set of academic talks will be organized during the EVA on the rooftop of Leggenda in Terzo, Italy has been organized. The first talk was conducted here in the video (link to the video on the front right).

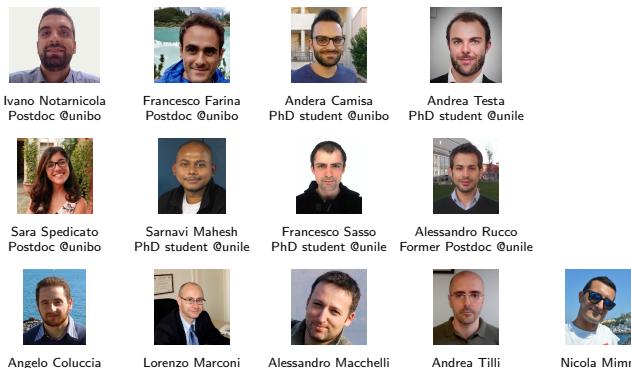
15 MAY 2018 Bologna
Brought to you by Prof. Daniel Zelený (Technion - Israel Institute of Technology, Israel). Specific problems from these four areas will be abstracted to a common mathematical set-up, and addressed by means of convex analysis, optimization methods, and graph theories. In particular, OPT4SMART will face the challenge of solving optimization problems under severe communication limitations, very-large-scale problem sizes, and non-convexity. The project will also promote the use of strong theoretical methods and effective numerical techniques available to people in engineering, Computer Science, Mathematics and other areas, who are facing optimization in cyber-physical

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Acknowledgment

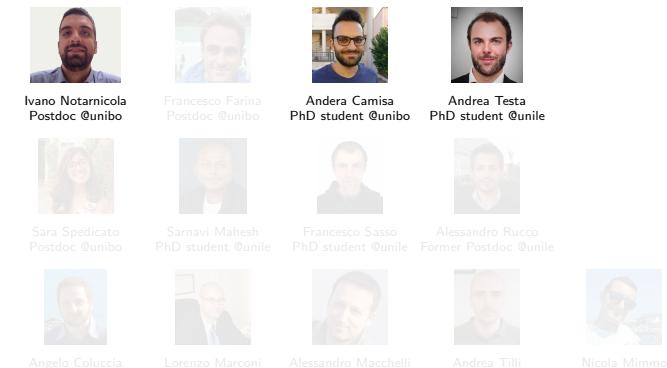


OPT4SMART group



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Acknowledgment



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Acknowledgment



Early collaborators / Mentors



F. Allgöwer



F. Bullo



M. Egerstedt



J. Hauser

Other collaborators

- F. Bayer, M. Bürger
- G. Scutari, Y. Sun
- R. Bouffanais, M. Chamanbaz
- M. Franceschelli
- A. Franchi
- N. Bof, R. Carli, L. Schenato, D. Varagnolo
- S. Chopra
- A. Garulli, A. Giannitrapani
- A. Falsone, M. Prandini
- K. Margellos, A. Papachristodoulou, L. Romao
- M. Bin, L. Marconi

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Take-home

Opportunities

- optimization is a building block in many estimation, learning, decision and control problems
- new powerful technology with massive computation and communication capability available

Challenges

- optimization (nonconvex, mixed-integer, combinatorial, big-data, stochastic, uncertain)
- network (asynchronous, unreliable, directed)



Take-home



Opportunities

- optimization is a building block in many estimation, learning, decision and control problems
- new powerful technology with massive computation and communication capability available

Distributed Optimization for Smart Cyber-Physical Networks

G. Notarstefano, I. Notarnicola, A. Camisa

Foundations and Trends® in Systems and Control (subm)

arxiv.org/abs/1906.10760

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